

# Euler Circuits in Graphs

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Here is why [3]:

- ▶ Denote each person as a vertex (node), and draw a line between two people if they are mutual acquaintances.
- ▶ Start from any vertex (labeled as  $u$ ), at least three vertices have edges (or no edges) with  $u$ , then we have a triangle (or empty triangle) if there are any edge between the three vertices, and they form an empty triangle (or triangle) if they have no edge at all.



# Seven Bridge Problem

The city of Königsberg in Prussia (now Kaliningrad, Russia) was set on both sides of the Pregel River, and included two large islands which were connected to each other and the mainland by seven bridges.

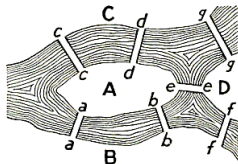


FIGURE 98. *Geographic Map:  
The Königsberg Bridges.*

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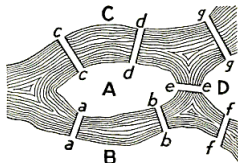
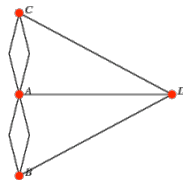


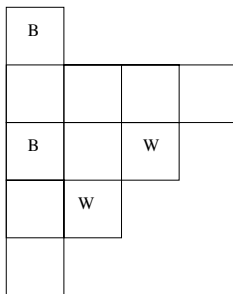
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## Knights on the chessboard

There are two black knights and two white knights on the partial chess board. Switch the black knights with the white knights (no two knights at same position any time). What is the minimum number of steps to do this?



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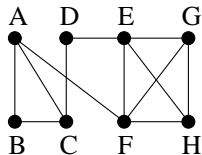
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- ▶ A graph  $G$  consists of vertices (nodes) and edges (lines) between vertices.
- ▶ In practice, vertices are objects (people, cities, et al), and edges are relations between objects (friendship, flights, et al).
- ▶ In a graph, every vertex  $x$  has neighbors—those vertices having an edge with  $x$ . We use  $N(x)$  to denote the neighbors of  $x$ . The **degree** of  $x$  is the number of edges incident to  $x$ .

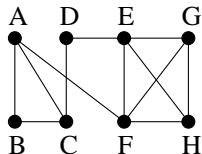
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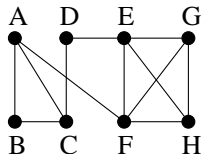


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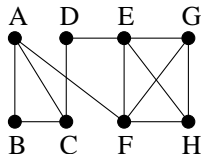
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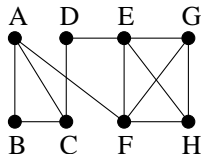
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- ▶ What are the degrees of the vertices in this graph?
- ▶ What is the meaning that  $A$  has degree three?
- ▶ For which houses, if any, is it possible to drive to all the other houses in less than 20 minutes?

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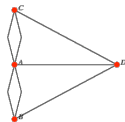
**Ex:** Can you draw a non-connected graph? What's the meaning of that?

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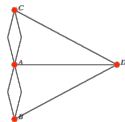
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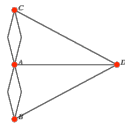
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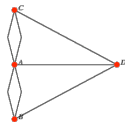
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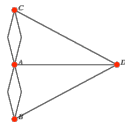
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- ▶ So if a graph is not connected, OR some vertex has an odd degree, then it has NO Euler circuit.



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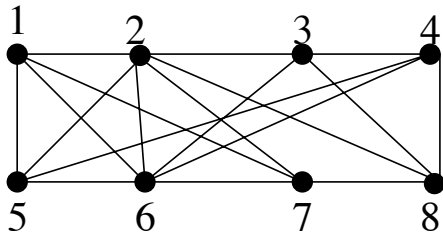
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  - ▶ Combine  $C$  with the Euler circuits to get an Euler circuit in the original graph

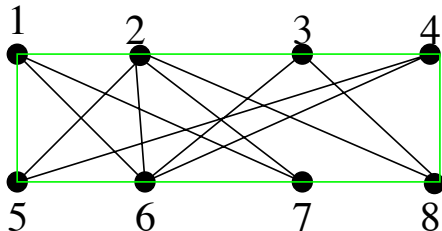
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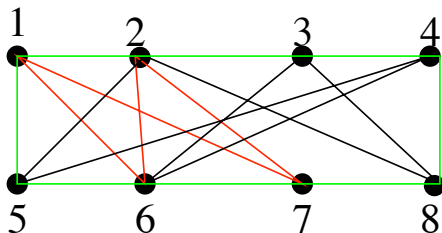
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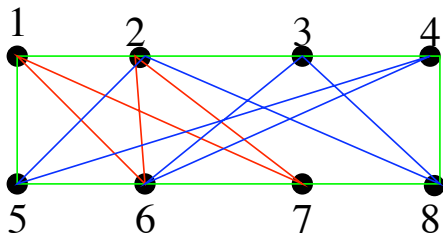


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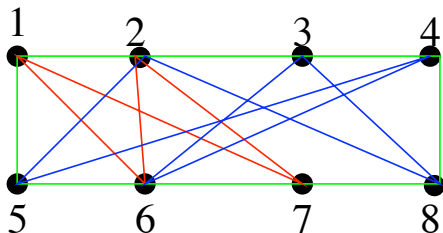
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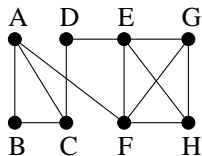
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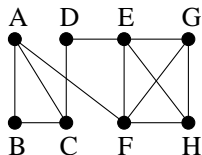
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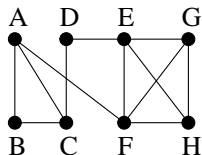
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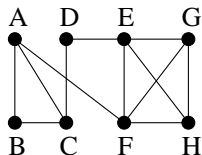
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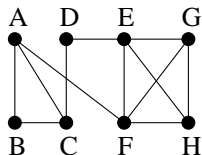
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- ▶ Observe that the vertices without even degrees are A, C, G, and H. Is this a coincidence?
- ▶ Not really! If add edges AC and GH to the original graph, we will get a graph with only even degrees, thus get an Euler circuit.

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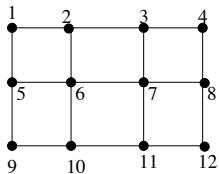
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- ▶ Yes! Here are two facts about graphs and the first one implies the second one
  - ▶ The sum of all degrees is even
  - ▶ The number of odd-degree vertices is even
- ▶ Now we know the number of odd-degree vertices is always even, so in order to create an Euler circuit in a non-Eulerian graph, we just need to **identify the odd-degree vertices, pair them up, add or repeat the corresponding edges.**

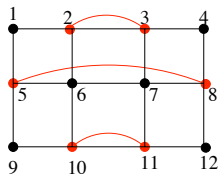
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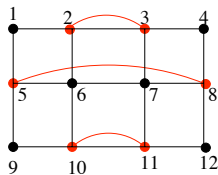
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- ▶ Identify the odd-degree vertices: 2, 3, 5, 8, 10, 12
- ▶ Pair up the vertices so that the total distance between pairs is smallest: (2, 3), (5, 8), and (10, 11) and add the edges

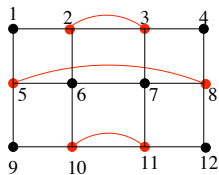


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- ▶ In the new graph, find an Euler circuit:  
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- ▶ Replace the added edges (red edges) with paths in the original graph:  
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- ▶ We will need **shortest path algorithm** and **minimum-weighted matching algorithm** to answer these two questions.